APRICOT 11:Multicast Workshop

Understanding & Deploying IP Multicast Networks

Srini Irigi, Stefan Olofsson, Greg Shepherd, SPG TME



Today's Agenda

- Multicast Fundamentals
- Multicast Service Models, Distribution Trees, Forwarding
- Multicast Protocol Basics
- Layer2 Multicast
- PIM Mechanics
- SSM
- MBGP
- MSDP

Fundamentals of IP Multicast



Agenda

- Why Multicast
- Multicast Applications
- Multicast Service Model
- Multicast Distribution Trees
- Multicast Forwarding
- Multicast Protocol Basics



Why Multicast?





Multicast Disadvantages

Multicast Is UDP Based!!!

- Best-effort delivery: Drops are to be expected. Multicast applications should not expect reliable delivery of data and should be designed accordingly. Reliable Multicast is still an area for much research. Expect to see more developments in this area.
- No congestion avoidance: Lack of TCP windowing and "slowstart" mechanisms can result in network congestion. If possible, Multicast applications should attempt to detect and avoid congestion conditions.
- Duplicates: Some multicast protocol mechanisms (e.g. Asserts, Registers and Shortest-Path Tree Transitions) result in the occasional generation of duplicate packets. Multicast applications should be designed to expect occasional duplicate packets.
- Out-of-sequence packets: Various network events can result in packets arriving out of sequence. Multicast applications should be designed to handle packets that arrive in some other sequence than they were sent by the source.

Multicast Service Model



IP Multicast Service Model

- RFC 1112 (Host Ext. for Multicast Support)
- Each multicast group identified by a class-D IP address
- Members of the group could be present anywhere in the Internet
- Members join and leave the group and indicate this to the routers
- Senders and receivers are distinct:

i.e., a sender need not be a member

 Routers listen to all multicast addresses and use multicast routing protocols to manage groups

IP Multicast Packet

Source address

Unique unicast IP address of the packet source

Destination address

ClassD address range

Does NOT represent a unique unicast destination address

Used to represent a unique group of receivers

IP Multicast Addressing

Multicast Group Addresses (224.0.0/4)

Range: 224.0.0.0–239.255.255.255

Old Class D address range.

High-order 4 bits are 1110

Multicast Address Ranges

- Link-Local Address Range 224.0.0.0–224.0.0.255
- Global Address Range
 224.0.1.0–238.255.255.255
- Administratively Scoped Address Range 239.0.0–239.255.255.25
- Scope Relative Address Range

Top 256 addresses of a Scoped Address Range

Link-Local Address Range

Assigned by IANA

224.0.0.0-224.0.0.255

Local wire multicast

TTL = 1

Examples:

224.0.0.5 = OSPF_DR's 224.0.0.10 = EIGRP Hello's 224.0.0.13 = All_PIM_Routers 224.0.0.22 = All_IGMPv3_Routers

Global Address Range

Assigned by IANA

Address Range: 224.0.1.0-238.255.255.255

Generally intended for "global" Internet scope multicast

Sometimes assigned to specific protocols

Example: Auto-RP (224.0.1.39 and 224.0.1.40)

Problem:

IANA is coming under increasing pressure from companies to assign them blocks of addresses for their applications or content services

This was never the intent of this block!

GLOP Addressing or SSM should be used instead!

Global Multicast Address Assignment

Dynamic Group Address Assignment

Historically accomplished using SDR application

Sessions announced over well-known group(s)

Address collisions detected and resolved at session creation time

Has problems scaling

Other techniques considered

Multicast Address Set-Claim (MASC)

Hierarchical, dynamic address allocation scheme

Unlikely to be deployed

No really good dynamic assignment method available for Global multicast

But is dynamic assignment really necessary with GLOP and SSM available?

Global Multicast Address Assignment

Static Group Address Assignment

RFC 3180—GLOP Addressing in 233/8

Group range: 233.0.0.0–233.255.255.255

Your AS number is inserted in middle two octets

Remaining low-order octet used for group assignment

EGLOP Addresses

Make use of private AS numbers

Assigned by a Registration Authority

Global Multicast Address Assignment

Static Group Address Assignment

Source Specific Multicast

Address range: 232.0.0.0/8

Flows based on both Group and Source address

Two different content flows can share the same Group address without interfering with each other

Provides virtually unlimited address space!

Preferred method for global one-to-many multicast

Private Multicast Address Assignment

 Assigned from the private 239.0.0.0/8 range May be subdivided into geographic scopes ranges Administration responsibility can be by scope range

Question:

"What technology is most often used to manage private multicast assignment?"

Answer:

A spreadsheet



Shortest Path or Source Tree



Shortest Path or Source Tree



Shared Tree



Shared Tree



Characteristics of Distribution Trees

Shortest Path trees

Uses more memory n(S x G) but you get optimal paths from source to all receivers; minimizes delay

Shared trees

Uses less memory n(G) but you may get sub-optimal paths from source to all receivers; may introduce extra delay

Multicast Forwarding



Unicast vs. Multicast Forwarding

Unicast Forwarding

Destination IP address directly indicates where to forward packet

Forwarding is hop-by-hop

Unicast routing table determines interface and next-hop router to forward packet

Unicast vs. Multicast Forwarding

Multicast Forwarding

Destination IP address (group) doesn't directly indicate where to forward packet

Forwarding is connection-oriented

Receivers must first be "connected" to the source before traffic begins to flow

Connection messages (PIM Joins) follow unicast routing table toward multicast source

Build Multicast Distribution Trees that determine where to forward packets

Distribution Trees rebuilt dynamically in case of network topology changes

The RPF Calculation

The multicast source address is checked against the unicast routing table

This determines the interface and upstream router in the direction of the source to which PIM Joins are sent

This interface becomes the "Incoming" or RPF interface

A router forwards a multicast datagram only if received on the RPF interface



RPF Calculation Based on Address of tree root Source or RP Best path to source found in **Unicast Route Table** Determines where to send Join Joins continue towards Source to build multicast tree Multicast data flows down tree Repeat for other receivers



RPF Calculation 10.1.1.1 SRC What if we have equal-cost paths? We can't use both **Tie-Breaker** Use highest Next-Hop IP address R 1.1.1.1 1.1.2.1 Join **Unicast Route Table** Ε2 Network Intfc Nxt-Hop 10.1.0.0/24 **E0** 1.1.1.1 **R1** 10.1.0.0/24 E1 1.1.2.1

Administrative Boundaries



Configured using the ip multicast boundary <acl>
 interface command

Administrative Boundaries



Multicast Protocol Basics



Types of Multicast Protocols

Dense-mode

Uses "Push" model

Traffic flooded throughout network

Pruned back where it is unwanted

Flood and prune behavior (typically every three minutes)

Sparse-mode

Uses "Pull" model

Traffic sent only to where it is requested

Explicit Join behavior

PIM-SM (RFC 4601)

Supports both source and shared trees

Assumes no hosts want multicast traffic unless they specifically ask for it

Uses a Rendezvous Point (RP)

Senders and Receivers "rendezvous" at this point to learn of each others existence

Senders are "registered" with RP by their first-hop router

Receivers are "joined" to the Shared Tree (rooted at the RP) by their local Designated Router (DR)

Appropriate for ...

Wide scale deployment for **both** densely and sparsely populated groups in the enterprise

Optimal choice for all production networks regardless of size and membership density
PIM-SM Shared Tree Join



PIM-SM Sender Registration



PIM-SM Sender Registration



PIM-SM Sender Registration















PIM-SM FFF

PIM-SM Frequently Forgotten Fact

"The default behavior of PIM-SM is that routers with directly connected members will join the Shortest Path Tree as soon as they detect a new multicast source."



PIM-SM—Evaluation

- Advantages:
 - Traffic only sent down "joined" branches
 - Can switch to optimal source-trees for high traffic sources dynamically
 - Unicast routing protocol-independent
 - Basis for inter-domain multicast routing
 - When used with MBGP and MSDP

Disadvantages

Few if any

Primary application

All production multicast networks with sparse or dense distribution of receivers

Protocol Summary

Conclusion

"Sparse mode good, dense mode bad!"

R. Davis "The Caveman's Guide to IP Multicast", 2000



IP Multicast at Layer 2



Module Agenda

- MAC Layer Multicast Addresses
- IGMPv2
- IGMPv3
- L2 Multicast Frame Switching IGMP Snooping PIM Snooping

MAC Layer Multicast Addresses



Layer 2 Multicast Addressing

IP Multicast MAC Address Mapping



Layer 2 Multicast Addressing IP Multicast MAC Address Mapping Be aware of the 32:1 address overlap

32—IP Multicast Addresses





IGMP

- How hosts tell routers about group membership
- Routers solicit group membership from directly connected hosts
- RFC 1112 specifies first version of IGMP
- RFC 2236 specifies IGMPv2

Most widely deployed and supported

RFC 3376 specifies IGMPv3

Growing support (required for SSM)

RFC 2236

Membership queries

Queries sent to 224.0.0.1 with ttl = 1

One router on LAN is elected to send queries

Query interval 60–120 seconds

Membership reports

IGMP report sent by one host suppresses sending by others

Restrict to one report per group per LAN

Unsolicited reports sent by host, when it first joins the group

RFC 2236

Group-specific query

Router sends Group-specific queries to make sure there are no members present before stopping to forward data for the group for that subnet

Leave Group message

Host sends leave message if it leaves the group and is the last member (reduces leave latency in comparison to v1)

RFC 2236

Querier election mechanism

On multi-access networks, an IGMP querier router is elected based on lowest IP address. Only the querier router sends queries.

Query-Interval Response Time

General queries specify "Max. Response Time" which inform hosts of the maximum time within which a host must respond to any query. (improves burstiness of the responses)

Backward compatible with IGMPv1

IGMPv2—Joining a Group



 Joining member sends report to 224.1.1.1 immediately upon joining (same as IGMPv1)

IGMPv2—Joining a Group



IGMPv2—Joining a Group



IGMPv2—Querier Election



- Initially all routers send out a query
- Router with lowest IP address "elected" querier
- Other routers become "non-queriers"

IGMPv2—Querier Election

Determining Which Router Is the IGMP Querier

rtr-a>show ip igmp interface e0 Ethernet0 is up, line protocol is up Internet address is 1.1.1.1, subnet mask is 255.255.255.0 IGMP is enabled on interface Current IGMP version is 2 CGMP is disabled on interface IGMP query interval is 60 seconds IGMP querier timeout is 120 seconds IGMP max query response time is 10 seconds Inbound IGMP access group is not set Multicast routing is enabled on interface Multicast TTL threshold is 0 Multicast designated router (DR) is 1.1.1.1 (this system) **IGMP querying router is 1.1.1.1 (this system)** Multicast groups joined: 224.0.1.40 224.2.127.254

IGMPv2—Maintaining a Group



- Router sends periodic queries
- One member per group per subnet reports
- Other members suppress reports





- H2 leaves group; sends Leave message
- Router sends Group specific query
- A remaining member host sends report
- Group remains active





- Last host leaves group; sends Leave message
- Router sends Group specific query
- No report is received
- Group times out





RFC 3376

Adds Include/Exclude Source Lists

Enables hosts to listen only to a specified subset of the hosts sending to the group

Requires new 'IPMulticastListen' API

New IGMPv3 stack required in the O/S

Apps must be rewritten to use IGMPv3 Include/Exclude features

Available in IOS 12.2, 12.1(3)T and 12.0(15)S
IGMPv3

RFC 3376

New membership report address

224.0.0.22 (All-IGMPv3-Routers)

All IGMPv3 hosts send reports to this address

Instead of the target group address as in IGMPv1/v2

All IGMPv3 routers listen to this address

Hosts do not listen or respond to this address

No report suppression

All hosts on wire respond to queries

Response Interval may be tuned over broad range

Useful when large numbers of hosts reside on subnet

IGMPv3 Example



H1—Member of 224.1.1.1

IGMPv3—Joining a Group



 Joining member sends IGMPv3 Report to 224.0.0.22 immediately upon joining

IGMPv3—Joining Specific Source(s)



- IGMPv3 report contains desired source(s) in the Include list
- Only "Included" source(s) are joined

IGMPv3—Excluding Specific Source(s)



- IGMPv3 report contains undesired source(s) in the Exclude list
- All sources except "Excluded" source(s) are joined

IGMPv3—Maintaining State



- Router sends periodic queries
- All IGMPv3 members respond

Reports contain multiple Group state records

PIM Sparse Mode



Module Agenda



- PIM Neighbor Discovery
- PIM State
- PIM SM Joining
- PIM SM Registering
- PIM SM SPT-Switchover

PIM Neighbor Discovery



PIM Neighbor Discovery



- PIMv2 Hellos are periodically multicast to the "All-PIM-Routers" (224.0.0.13) group address (default = 30 seconds)
- If the "DR" times-out, a new "DR" is elected
- The "DR" is responsible for sending all Joins and Register messages for any receivers or senders on the network

PIM Neighbor Discovery—IOS

wan-gw8> show ip pim neighbor PIM Neighbor Table				
Neighbor	Interface	Uptime/Expires	Ver	Mode
Address				Prio/Mode
171.68.0.70	FastEthernet0/0	2w1d/00:01:24	v2	1 / B S
171.68.0.91	FastEthernet0/0	2w6d/00:01:01	v2	1 / B S
171.68.0.82	FastEthernet0/0	7w0d/00:01:14	v2	5 / DR B S
171.68.0.86	FastEthernet0/0	7w0d/00:01:13	v2	1 / B S
171.68.0.80	FastEthernet0/0	7w0d/00:01:02	v2	1 / B S
171.68.28.70	Serial2.31	22:47:11/00:01:16	v2	1 / B S
171.68.28.50	Serial2.33	22:47:22/00:01:08	v2	1 / B S
171.68.27.74	Serial2.36	22:47:07/00:01:21	v2	N /
171.68.28.170	Serial0.70	1d4h/00:01:06	v2	N /
171.68.27.2	Serial1.51	1w4d/00:01:25	v2	1 / B S
171.68.28.110	Serial3.56	1d4h/00:01:20	v2	1 / B S
171.68.28.58	Serial3.102	12:53:25/00:01:03	v2	1 / B S

DR Failover



Expiration time sent in PIM query messages

Expiration time = 3 x <query-interval>

Default <query-interval> = 30 seconds

DR failover ~ 90 seconds (worst case) by default

Tuning DR Failover

Tune PIM query interval

Use interface configuration command

ip pim query-interval <period> [msec]

Default <period> = seconds

"msec" keyword available beginning with 12.1(11b)E

Permits DR failover to be adjusted

Sub-second DR failover possible

Smaller intervals increase PIM query traffic

Increase is usually insignificant

PIM State



PIM State

- Describes the "state" of the multicast distribution trees as understood by the router at this point in the network
- Represented by entries in the multicast routing (mroute) table
 - Used to make multicast traffic forwarding decisions
 - Composed of (*, G) and (S, G) entries
 - Each entry contains RPF information
 - Incoming (i.e. RPF) interface
 - RPF Neighbor (upstream)
 - Each entry contains an Outgoing Interface List (OIL)
 - OIL may be NULL

PIM-SM State Example—IOS

sj-mbone> show ip mroute Flags: D - Dense, S - Sparse, B - Bidir Group, s - SSM Group, C - Connected, L - Local, P - Pruned, R - RP-bit set, F - Register flag, T - SPT-bit set, J - Join SPT, M - MSDP created entry, X - Proxy Join Timer Running, A - Candidate for MSDP Advertisement, U - URD, I - Received Source Specific Host Report Outgoing interface flags: H - Hardware switched Timers: Uptime/Expires Interface state: Interface, Next-Hop or VCD, State/Mode (*, 224.1.1.1), 2w1d/00:00:00, RP 172.16.25.1, flags: SJC Incoming interface: Serial0/1, RPF nbr 172.16.4.1 Outgoing interface list: Ethernet0/1, Forward/Sparse-Dense, 2w1d/00:01:40 Serial0/0, Forward/Sparse-Dense, 00:4:52/00:02:08 (172.16.8.2, 224.1.1.1), 00:00:10/00:02:59, flags: CJT Incoming interface: Serial0/1, RPF nbr 172.16.4.1 Outgoing interface list: Ethernet0/1, Forward/Sparse-Dense, 00:00:10/00:02:49 Serial0/0, Forward/Sparse-Dense, 00:4:52/00:02:08

PIM-SM (*,G) State Rules

(*,G) creation

Receipt of a (*,G) Join or IGMP Report Automatically if (S,G) must be created

(*,G) reflects default group forwarding

IIF = RPF interface toward RP

OIL = interfaces

That received a (*,G) Join or

With directly connected members or

Manually configured

(*,G) deletion

When OIL = NULL and

No child (S,G) state exists

PIM-SM (S,G) State Rules

(S,G) creation

By receipt of (S,G) Join or Prune or

By "Register" process

Parent (*,G) created (if doesn't exist)

 (S,G) reflects forwarding of "S" to "G" IIF = RPF Interface normally toward source RPF toward RP if "RP-bit" set OIL = Initially, copy of (*,G) OIL minus IIF
(S,G) deletion

By normal (S,G) entry timeout

PIM-SM OIL Rules

Interfaces in OIL added

By receipt of Join message

Interfaces added to (*,G) are added to all (S,G)s

Interfaces in OIL removed

By receipt of Prune message

Interfaces removed from (*,G) are removed from all (S,G)s

Interface expire timer counts down to zero

Timer reset (to 3 min.) by receipt of periodic Join

or

By IGMP membership report

PIM-SM Triggered Join/Prune Rules

Triggering Join/Prune Messages

(*,G) Joins are triggered when:

The (*,G) OIL transitions from Null to non-Null

(*,G) Prunes are triggered when:

The (*,G) OIL transitions from non-Null to Null

(S,G) Joins are triggered when:

The (S,G) OIL transitions from Null to non-Null

(S,G) Prunes are triggered when:

The (S,G) OIL transitions from non-Null to Null

(S,G)RP-bit Prunes are triggered when:

The (S,G) RPF info != the (*,G) RPF info

PIM-SM State Flags

- S = Sparse
- C = Directly Connected Host
- L = Local (Router is member)
- P = Pruned (All intfcs in OIL = Prune)
- T = Forwarding via SPT

Indicates at least one packet was forwarded

PIM-SM State Flags (Cont.)

J = Join SPT

In (*, G) entry

Indicates SPT-Threshold is being exceeded

Next (S,G) received will trigger join of SPT

In (S, G) entry

Indicates SPT joined due to SPT-Threshold

If rate < SPT-Threshold, switch back to Shared Tree

F = Register/First-Hop

In (S,G) entry

"S" is a directly connected source

Triggers the Register Process

In (*, G) entry

Set when "F" set in at least one child (S,G)

PIM-SM State Flags (Cont.)

R = RP bit

(S, G) entries only

Set by (S,G)RP-bit Prune

Indicates info is applicable to Shared Tree

Used to prune (S,G) traffic from Shared Tree

Initiated by Last-hop router after switch to SPT

Modifies (S,G) forwarding behavior

IIF = RPF toward RP (I.e. up the Shared Tree)

OIL = Pruned accordingly

PIM SM Joining





① Rcvr wishes to receive group G traffic. Sends IGMP Join for G.





① Rcvr wishes to receive group G traffic. Sends IGMP Join for G.

② B sends (*,G) Join towards RP.



A Creates (*, 224.1.1.1) State



Output to the second second

- B sends (*,G) Join towards RP.
- ③ A sends (*,G) Join towards RP.
- ④ Shared tree is built all the way back to the RP.

PIM SM Registering



PIM SM Register Scenarios

- Receivers Join Group First
- Source Registers First
- Receivers along the SPT





State in "RP" Before Any Source Registers (With Receivers on Shared Tree)



State in B Before Any Source Registers (With Receivers on Shared Tree)



State in A Before Any Source Registers (With Receivers on Shared Tree)



① Source begins sending group G traffic.


A Creates (S, G) State for Source

(After Automatically Creating a (*, G) entry)

Source begins sending group G traffic.

2 A encapsulates packets in Registers; unicasts to RP.



"RP" Processes Register; Creates (S, G) State

RP (C) de-encapsulates packets; forwards down Shared tree.





TECRST-1008_c1 © 2009 Cisco Systems, Inc. All rights reserved. Cisco Public



B sends (S,G) Join toward Source to continue building SPT.



A Processes the (S, G) Join; Adds Serial0 to OIL



O RP begins receiving (S,G) traffic down SPT.



Traffic Arriving via SPT Is Forwarded Down Shared Tree (This Causes the "T" Flag to Be Set)



Once "T" Flag is set, next "Register" causes RP to send back a "Register-Stop" to A



A Stops Sending Register Messages (Final State in A)

(S,G) Traffic now flowing down a single path (SPT) to RP.



Final State in B



Final State in the "RP" (With Receivers on Shared Tree)





(Without Receivers on Shared Tree)



State in B Before Any Source Registers (With Receivers on Shared Tree)



State in A Before Any Source Registers (With Receivers on Shared Tree)



(1) Source begins sending group G traffic.



A Creates (S, G) State for Source

(After Automatically Creating a (*, G) Entry)
Source begins sending group G traffic.
A encapsulates packets in Registers; unicasts to RP.



③ RP (C) has no receivers on Shared Tree; discards packet.



④ RP sends "Register-Stop" to A.



A stops encapsulating traffic in Register Messages; drops packets from Source.



State in A After Registering (Without Receivers on Shared Tree)



State in B After A Registers (Without Receivers on Shared Tree)



State in RP After A Registers (Without Receivers on Shared Tree)





Receivers Begin Joining the Shared Tree





RP Processes (*,G) Join (Adds Serial1 to Outgoing Interface Lists)

RP sends (S,G) Joins for all known Sources in Group.



B Processes Join, Creates (S, G) State (After Automatically Creating the (*, G) Entry)

B sends (S,G) Join toward Source to continue building SPT.



A Processes the (S, G) Join; Adds Serial0 to OIL





Final State in Router A (IOS)



Final State in B After Receivers Join



Final State in RP After Receivers Join (IOS)





Current State in B



Current State in the RP



① Rcvr wishes to receive group G traffic. Sends IGMP Join for G.



State in B After Rcvr Joins Group




PIM SM Registering Receivers Along the SPT



State in RP After B Joins Shared Tree

PIM SM Registering Receivers Along the SPT



③ Group G traffic begins to flow to Rcvr.

(Note: 171.68.37.121 traffic doesn't flow to RP then back down to B)



- SPT Thresholds may be set for any Group Access Lists may be used to specify which Groups Default Threshold = 0kbps (I.e. immediately join SPT) Threshold = "infinity" means "never join SPT"
 Don't use values in between "0" and "infinity" (In IOS XR, "0" and "infinity" are the only options)
- Threshold triggers Join of Source Tree
 Sends an (S,G) Join up SPT for next "S" in "G" packet received



State in C Before Switch



State in D Before Switch



State in A Before Switch



State in B Before Switch



Mew source (S_i,G) packet arrives down Shared tree.



B creates (S_i,G) state.

TECRST-1008_c1 © 2009 Cisco Systems, Inc. All rights reserved. Cisco Public



B sends (S_i,G) Join towards S_i.





(4) A triggers (S_i ,G) Join toward S_i .



A triggers (S_i,G) Join toward S_i. (S_i, G) traffic begins flowing down SPT tree.





Once "T" flag is set, A triggers (S_i,G)RP-bit Prunes toward RP.





Unnecessary (S_i, G) traffic is pruned from the Shared tree.

PIM SM Pruning





State in B Before Pruning



State in A Before Pruning



Is a Leaf router. Last Rcvr, leaves group G.





B removes E1 from (*,G) and any (S_i,G) "oilists".

B's (*,G) "oilist" now empty; triggers (*,G) Prune toward RP.





 A receives Prune; removes E0 from (*,G) "oilist". (After the 3 second Multi-access Network Prune delay.)
 A's (*,G) "oilist" now empty; triggers (*,G) Prune toward RP.







State in A Before Pruning



State in B Before Pruning



Is a Leaf router. Last Rcvr leaves group G.





B removes E1 from (*,G) and all (S,G) OILs.



B's (*,G) OIL now empty; triggers (*,G) Prune toward RP.



B's (S,G) OIL also now empty; triggers (S, G) Prune towards Si.



A receives (*, G) Prune; removes E0 from (*,G) & (S,G) OILs (After the 3 second Multi-access Network Prune delay.)



(6) A' s (*,G) *OIL* now empty; triggers (*,G) Prune toward RP.


 \bigcirc A's (S,G) OIL also now empty; triggers (S,G) Prune towards S_i.



(S_i,G) traffic ceases flowing down SPT.

Recap: Common Multicast Flags— IOS

- S: Sparse Mode (in contrast to D for Dense Mode)
- s: SSM; only seen on (S,G) entries
- B: Bidir
- F: Register; set on first-hop router
- P: Prune; entry has an empty OIL
- J: Join-SPT; (*,G) traffic exceeds SPT Threshold
- T: SPT; set on (S,G) entries after SPT join
- L: Local; router should receive and process this traffic
- C: Connected; seen primarily with IGMP

Source Specific Multicast



Barriers to Multicast Deployment

Global Multicast Address Allocation

Dynamic Address Allocation

No adequate dynamic address allocation methods exist

SDR—Doesn't scale

MASC—Long ways off!

Static Address Allocation (GLOP)

Based on AS number

Insufficient address space for large Content Providers

Multicast Content "Jammers"

Undesirable sources on a multicast group

"Capt. Midnight" sources bogus data/noise to group

Can cause DoS attack by congesting low speed links

Source Specific Multicast (SSM)

- Uses Source Trees only
- Assumes one-to-many model Most Internet multicast fits this model IP/TV also fits this model
- Hosts responsible for source discovery Typically via some out-of-band mechanism Web page, Content Server, etc.
 Eliminates need for RP and Shared Trees Eliminates need for MSDP

SSM Overview

- Hosts join a specific source within a group Content identified by specific (S,G) instead of (*,G) Hosts responsible for learning (S,G) information
- Last-hop router sends (S,G) join toward source Shared Tree is never Joined or used Eliminates possibility of content Jammers Only specified (S,G) flow is delivered to host
- Eliminates Networked-Based Source Discovery

No RPs for SSM groups

Simplifies address allocation

Dissimilar content sources can use same group without fear of interfering with each other

SSM Example



SSM Example



SSM Configuration

Global command

ip pim ssm {default | range <acl>}

Defines SSM address range

Default range = 232.0.0.0/8

Use ACL for other ranges

Prevents Shared Tree Creation

(*, G) Joins never sent or processed

PIM Registers never sent or processed

Available in Cisco IOS versions

12.1(5)T, 12.2, 12.0(15)S, 12.1(8)E

SSM—Summary

Uses Source Trees only

Hosts are responsible for source and group discovery Hosts must signal router which (S,G) to join

Solves multicast address allocation problems

Flows differentiated by **both** source and group

Content providers can use same group ranges

Since each (S,G) flow is unique

- Helps prevent certain DoS attacks
 - "Bogus" source traffic:

Can't consume network bandwidth

Not received by host application

Bidirectional (BiDir) PIM



Multicast Application Categories

One-to-many applications

Video, TV, radio, concerts, stock ticker, etc.

Few-to-few applications

Small (<10 member) video/audio conferences

- Few-to-many applications
 TIBCO RV servers (publishing)
- Many-to-many applications
 Stock trading floors, gaming
- Many-to-few applications
 TIBCO RV clients (subscriptions)

Multicast Application Categories PIM-SM (S, G) State

- One-to-many applications
 Single (S,G) entry
- Few-to-few applications
 Few (<10 typical) (S,G) entries
- Few-to-many applications
 Few (<10 typical) (S,G) entries
- Many-to-many applications
 Unlimited (S,G) entries
- Many-to-few applications
 Unlimited (S,G) entries

Many-to-Any State Problem

Creates huge amounts of (S,G) state State maintenance workloads skyrocket High OIL fan-outs make the problem worse Router performance begins to suffer Using Shared-Trees only Provides some (S,G) state reduction Results in (S,G) state only along SPT to RP Frequently still too much (S,G) state Need a solution that only uses (*,G) state

Bidirectional (BiDir) PIM

Idea:

Use the same tree for traffic from sources towards RP and from RP to receivers

Benefits:

Less state in routers

Only (*, G) state is used

Source traffic follows the Shared Tree

Flows up the Shared Tree to reach the RP

Flows down the Shared Tree to reach all other receivers

Bidirectional (BiDir) PIM

Bidirectional Shared-Trees
 Violates current (*,G) RPF rules
 Traffic often accepted on outgoing interfaces
 Care must be taken to avoid multicast loops
 Requires a Designated Forwarder (DF)
 Responsible for forwarding traffic up Shared Tree
 DFs will accept data on the interfaces in their OIL
 Then send it out all other interfaces (including the IIF)

Bidirectional PIM—Overview



Bidirectional PIM—Overview



PIM Modifications for BiDir Operation

Designated Forwarders (DF)

One DF per link

Router with best path to the RP is elected DF

Note: Designated Routers (DR) are not used for bidir groups

In addition to normal (*,G) forwarding rules:

Accepts traffic on outgoing interfaces

Forwards traffic out all other interfaces

Designated Forwarder Election

- Automatically performed on every link
 When Bidir Group-range/RP is learned or configured
 Router with the best path to the RP elected DF
 Uses assert-like metric comparison to pick best path

 Purpose:
 - Ensures all routers on link agree on who is DF Prevents route loops from forming



Receiver 1 Joins Group Causing Router "D" to Create (*, G) State



Router "D" Sends (*, G) Join to Router "F" (DF) Causing It to Create (*, G) State



Router "F" Sends (*, G) Join to "RP" Causing It to Create (*, G) State



Branch of Shared Tree Is Now Built Down to Receiver 1



Receiver 2 Also Joins Group



Router "B" Creates (*, G) State



Router "B" Sends (*, G) Join to "E" (DF) Causing It to Create (*, G) State



Router "E" Sends (*, G) Join to "RP" (State on RP Remains Unchanged)



New Branch of Shared Tree Is Built to Receiver 2



Arriving Traffic from Source Causes Router "A" to Create (*, G) State



Traffic Is Forwarded Toward Router "E" and Also Arrives at IIF of Router "B"



Router "B" Forwards Traffic Back Down Shared Tree ala Normal PIM-SM



Router "E" Forwards Traffic on Toward RP



Traffic Forwarded Toward RP Also Arrives at the IIF of Router "F"



Router "F" Forwards Traffic on Down the Shared Tree ala Normal PIM-SM
Forwarding/Tree Building



Router "D" Forwards Traffic to Receiver 1 via the Shared Tree

Forwarding/Tree Building



Question: Does the RP even have to physically exist?

Forwarding/Tree Building



Question: Does the RP even have to physically exist? Answer: No. It can just be a phantom address.

Bidir PIM—Summary

Uses Shared Trees only

Single (*, G) forwarding entry per group Source traffic flows up and down Shared Tree

Drastically reduces network mroute state
 Eliminates ALL (S,G) state in the network
 By eliminating SPT between source and RP
 Allows many-to-any applications to scale
 Permits virtually an unlimited number of sources

Lab



IPv4 PIM Configuration

Enable multicast routing on every router

ip multicast-routing

ALL Modes of PIM on ALL interfaces of every router

interface <interface>

ip pim sparse-mode

Sparse-mode and BiDir require an RP mapping on every router

ip pim rp-address x.x.x.x [bidir]

SSM: 232/8 is the default range

ip pim ssm default

On the RP router ONLY

interface lo1

ip address x.x.x.x 255.255.255.255

LAB #1 PIM-SM Mechanics - SSM / ASM / BiDir

- Get your username and password from the instructor
- Once your are logged in, DO NOT start the lab until instructed
- Lab templates PIM-Mechanics
- Refer to your lab handout

LAB #1 IPv4 PIM-SM Mechanics - SSM / ASM / BiDir



IPv4 vs. IPv6 Multicast

IP Service	IPv4 Solution	IPv6 Solution	
Address Range	32-Bit, Class D	128-Bit (112-Bit Group)	
	Protocol-Independent	Protocol-Independent	
Routing	All IGPs and GBP4+	All IGPs and BGP4+ with v6 Mcast SAFI	
Forwarding	PIM-DM, PIM-SM: ASM, SSM, BiDir	PIM-SM: ASM, SSM, BiDir	
Group Management	IBMPv1, v2, v3	MLDv1, v2	
Domain Control	Boundary/Border	Scope Identifier	
Interdomain Source Discovery	MSDP Across Independent PIM Domains	Single RP Within Globally Shared Domains	

IPv6 Multicast Addresses (RFC 3513)



IPv6 Layer 2 Multicast Addressing Mapping



Unicast-Based Multicast Addresses

8	4	4	8	8	64	32
FF	Flags	Scope	Rsvd	Plen	Network-Prefix	Group-ID

RFC 3306—unicast-based multicast addresses

Similar to IPv4 GLOP addressing

Solves IPv6 global address allocation problem

Flags = 00PT

 $P = 1, T = 1 \rightarrow$ Unicast-based multicast address

Example

Content provider's unicast prefix

1234:5678:9::/48

Multicast address

FF3x:0030:1234:5678:0009::0001

IP Routing for Multicast

- RPF-based on reachability to v6 source same as with v4 multicast
- RPF still protocol-independent

Static routes, mroutes

Unicast RIB: BGP, ISIS, OSPF, EIGRP, RIP, etc.

Multiprotocol BGP (mBGP)

Support for v6 mcast subaddress family

Provide translate function for nonsupporting peers

IPv6 Multicast Forwarding

- PIM-Sparse Mode (PIM-SM)
 RFC4601
- PIM Source Specific Mode (SSM)

RFC3569 SSM overview (v6 SSM needs MLDv2)

Unicast, prefix-based multicast addresses ff30::/12

SSM range is ff3X::/96

PIM Bi-Directional Mode (BiDir)

draft-ietf-pim-bidir-09.txt

RP Mapping Mechanisms for IPv6

- Static RP assignment
- BSR
- Auto-RP—no current plans
- Embedded RP

Embedded RP Addressing—RFC3956

8	4	4	4	4	8	64	32
FF	Flags	Scope	Rsvd	RPadr	Plen	Network-Prefix	Group-ID

Proposed new multicast address type

Uses unicast-based multicast addresses (RFC 3306)

- RP address is embedded in multicast address
- Flag bits = 0RPT

R = 1, P = 1, T = 1 \rightarrow Embedded RP address

- Network-Prefix::RPadr = RP address
- For each unicast prefix you own, you now also own:

16 RPs for each of the 16 multicast scopes (256 total) with 2^32 multicast groups assigned to each RP (2^40 total)

Embedded RP Addressing—Example

Multicast Address with Embedded RP Address



Multicast Listener Discover—MLD

- MLD is equivalent to IGMP in IPv4
- MLD messages are transported over ICMPv6
- Version number confusion

MLDv1 corresponds to IGMPv2

RFC 2710

MLDv2 corresponds to IGMPv3, needed for SSM

RFC 3810

MLD snooping

draft-ietf-magma-snoop-12.txt

IPv6 PIM Configuration

Enable multicast routing on every router

ipv6 multicast-routing

ALL Modes of PIM on ALL interfaces of every router

interface <interface>

ip pim sparse-mode

Sparse-mode and BiDir require an RP mapping on every router

ipv6 pim rp-address x.x.x.x [bidir]

On the RP router ONLY

interface lo1

ipv6 address XXX::XXX/128

LAB #1 IPv6 PIM-SM Mechanics - SSM / ASM / BiDir



Agenda

- Introduction
- Multicast addressing
- Group Membership Protocol
- PIM-SM / SSM
- MSDP
- MBGP
- Summary

Uses inter-domain source trees only.

RP's know about all sources in their domain Sources cause a "PIM Register" to the RP Can tell RP's in other domains of its sources Via MSDP SA (Source Active) messages

RP's know about receivers in their domain Receivers cause a "(*, G) Join" to the RP RP can join the source tree in the peer domain Via normal PIM (S, G) joins Only necessary if there are receivers for the group Last-hop routers then join source tree directly.



MSDP Example



Domain E



MSDP Example Domain E RP **MSDP** Peers **Multicast Traffic** Domain C RP Domain B RP RP Domain D RP S Domain A



MSDP Example Domain E RP • **MSDP** Peers Multicast Traffic Domain C RP Domain B RP RP Domain D RP S Domain A

- MSDP Peers configured similar to BGP
- Peers connect using TCP port 639

Lower address peer initiates connection

Higher address peer waits in LISTEN state

- Peers send keepalives every 60 secs.
- Connection reset after 75 seconds

If no MSDP packets or keepalives are received

• MSDP peers normally must run BGP!

BGP NLRI is used to RPF check SA messages.

May use NLRI from M-Table, U-Table or both.

RPF check prevents SA's from looping.

(More on that later.)

Exceptions:

When peering with only a single MSDP peer. When using an MSDP Mesh-Group.



 MSDP peer connections are established using the MSDP "peer" configuration command

ip msdp peer <ip-address> [connect-source <intfc>]



 MSDP peer connections are established using the MSDP "peer" configuration command

ip msdp peer <ip-address> [connect-source <intfc>]



 MSDP peer connections are established using the MSDP "peer" configuration command

ip msdp peer <ip-address> [connect-source <intfc>]



• Stub-networks may use "default" peering without being a BGP peer by using the MSDP "default-peer" configuration command.

ip msdp default-peer <ip-address>



• Multiple "default-peers" may be configured in case connection to first default-peer goes down.



• When connection to first 'default-peer' is lost, the next one in the list is tried.


 Stub-networks configured with only a single MSDP peer are treated in the same manner as when a single "default-peer" is configured. (i.e. BGP is not required.)

SA Message Contents

MSDP Source Active (SA) Messages
 Used to advertise active Sources in a domain
 Can also carry 1st multicast packet from source
 Hack for Bursty Sources (a' la SDR)

 SA Message Contents:
 IP Address of Originating RP
 Number of (S, G)' s pairs being advertised
 List of active (S, G)' s in the domain
 Encapsulated Multicast packet [optional]

Local Sources

RP's only originate SA's for local sources Denoted by the "A" flag on an (S,G) entry on RP A source is local if:

The RP received a "Register" for (S, G), or

The source is directly connected to RP

Use 'msdp redistribute' to control what SA's are originated.

Think of this as 'msdp sa-originate-filter' function

ip msdp redistribute [list <acl>]
 [asn <aspath-acl>]
 [route-map <map>]

Filter by (S,G) pair using 'list <acl>'

Filter by AS-PATH using 'asn <aspath-acl>'

Filter based on route-map '<map>'

Omitting all acl's stops all SA origination

Example: ip msdp redistribute

Default: Originate SA's for all local sources

If 'msdp redistribute' command is not configured

 SA messages are triggered when any new source in the local domain goes active.

Initial multicast packet is encapsulated in an SA message.

This is an attempt at solving the bursty-source problem

Encapsulating Initial Multicast Packets

Can bypass TTL-Thresholds

Original TTL is inside of data portion of SA message

SA messages sent via Unicast with TTL = 255

Requires special command to control

ip msdp ttl-threshold <peer-address> <ttl>

Encapsulated multicast packets with a TTL lower than <ttl> for the specific MSDP peer are not forwarded or originated.

Once a minute

Router scans mroute table

If group = sparse AND router = RP for group For each (S,G) entry for the group: If the 'msdp redistribute' filters permits AND if the source is a local source Then originate an SA message for (S,G)

Receiving SA Messages

If SA message RPF checks OK Store in SA Cache If new SA cache entry Immediately flood SA downstream Set entry's SA-expire-timer to 6 minutes. If RP for group and receivers exist Create (S,G) entry and trigger (S,G) Join If existing entry **Reset entry's SA-expire-timer to 6 minutes.** When timer = zero, entry has expired and is deleted.

Else

Discard SA

SA Message Cache

Enabling SA Caching

ip msdp cache-sa-state [list <acl>]

Caching is now on by default.

Beginning with IOS versions 12.1(7), 12.0(14)S1.

Cannot be turned off.

Router caches all SA messages.

Cached (S, G) entries timeout after 6 minutes.

If not refreshed by another (S,G) SA message.

Once per minute, router scans SA cache.

Sends SA downstream for each entry in cache.

SA Message Caching

Listing the contents of the SA Cache

show ip msdp sa-cache [<group-or-source>] [<asn>]

sj-mbone# show ip msdp sa-cache
MSDP Source-Active Cache - 1997 entries
(193.92.8.77, 224.2.232.0), RP 194.177.210.41, MBGP/AS 5408, 00:01:51/00:04:09
(128.119.167.221, 224.77.0.0), RP 128.119.3.241, MBGP/AS 1249, 06:40:59/00:05:12
(147.228.44.30, 233.0.0.1), RP 195.178.64.113, MBGP/AS 2852, 00:04:48/00:01:11
(128.117.16.142, 233.0.0.1), RP 204.147.128.141, MBGP/AS 145, 00:00:41/00:05:18
(132.250.95.60, 224.253.0.1), RP 138.18.100.1, MBGP/AS 668, 01:15:07/00:05:55
(128.119.40.229, 224.2.0.1), RP 128.119.3.241, MBGP/AS 1249, 06:40:59/00:05:12
(130.225.245.71, 227.37.32.1), RP 130.225.245.71, MBGP/AS 1835, 1d00h/00:05:29
(194.177.210.41, 227.37.32.1), RP 194.177.210.41, MBGP/AS 5408, 00:02:53/00:03:07
(206.190.42.106, 236.195.60.2), RP 206.190.40.61, MBGP/AS 5779, 00:07:27/00:04:04

Clearing the contents of the SA Cache

clear ip msdp sa-cache [<group-address> | group-name]

Filtering Incoming/Outgoing SA Messages

SA Filter Command:

```
ip msdp sa-filter {in|out} <peer-address> [list <acl>]
        [route-map <map>]
```

Filters (S,G) pairs to / from peer based on specified ACL.

Can filter based on AS-Path by using optional route-map clause with a path-list acl.

You can filter flooded and originated SA's based on a specific peer, incoming and outgoing.

Caution: Filtering SA messages can break the Flood and Join mechanism!

Recommended MSDP SA Filter



See "ftp://ftp-eng.cisco.com/ipmulticast/msdp-safilter.txt" for the latest updates to this list.

SA Message RPF Checking

Purpose

Accept SA's via a single deterministic path

Ignore all other arriving SA's

Necessary to prevent SA's from looping endlessly

Problem

Need to know MSDP topology of Internet

But, MSDP does not distribute topology data!

Solution

Use BGP data to *infer* MSDP topology.

Impact:

The MSDP topology must follow BGP topology.

TECRST-1008_c1 © 2009 Cisco Systems Achi MSPP neer 265

SA Message RPF Checking

RPF Check Rules depend on peering

Rule 1: Sending MSDP peer = iBGP peer

Rule 2: Sending MSDP peer = eBGP peer

Rule 3: Sending MSDP peer != BGP peer

Exceptions:

RPF check is skipped when: Sending MSDP peer = Originating RP Sending MSDP peer = Mesh-Group peer Sending MSDP peer = only MSDP peer (i.e. the 'default-peer' or the only 'msdp-peer' configured.)

SA Message RPF Checking

Determining Applicable RPF Rule

Use IP address of sending MSDP peer

Find BGP neighbor w/matching IP address

IF (no match found)

Apply Rule 3

IF (matching neighbor = iBGP peer)

Apply Rule 1

ELSE {matching neighbor = eBGP peer}

Apply Rule 2

Implication

The MSDP peer address must be configured using the same IP address as the BGP peer!

•When MSDP peer = GBGP peer

Find "Best Path" to RP in BGP Tables

Search M-Table first then U-Table.

If no path to Originating RP found, RPF Fails

Note "BGP Neighbor" that advertised path

(i.e IP Address of BGP peer that sent us this path)

Warning:

This is not the same as the Next-hop of the path!!!

iBGP peers normally do not set Next-hop = Self.

This is also not necessarily the same as the Router-ID!

Rule 1 Test Condition:

MSDP Peer address = BGP Neighbor address?

recRST-1008_c1 © 2009 Cisco Systelif, InY es, reRPF Succeeds



TECRST-1008_c1 © 2009 Cisco Systems, Inc. All rights reserved.



TECRST-1008_c1 © 2009 Cisco Systems, Inc. All rights reserved.





RPF Check Rule 2

When MSDP peer = eBGP peer

Find BGP "Best Path" to RP

Search M-Table first then U-Table.

If no path to Originating RP found, RPF Fails

Rule 2 Test Condition:

First AS in path to the RP = AS of eBGP peer?

If Yes, RPF Succeeds







RPF Check Rule 3 When MSDP peer != BGP peer

Find BGP "Best Path" to RP

Search M-Table first then U-Table.

If no path to Originating RP found, RPF Fails

Find BGP "Best Path" to MSDP peer

Search M-Table first then U-Table.

If no path to sending MSDP Peer found, RPF Fails

Note AS of sending MSDP Peer

Origin AS (last AS) in AS-PATH to MSDP Peer

Rule 3 Test Condition:

First AS in path to RP = Sending MSDP Peer AS ?

If Yes, RPF Succeeds





MSDP Mesh-Groups

Optimises SA flooding.

Useful when 2 or more peers are in a group.

Requires full mesh of mesh group peers.

Reduces amount of SA traffic in the net.

SA's not flooded to other mesh-group peers.

Suspends RPF check of SA messages.

When received from a mesh-group peer.

SA's always accepted from mesh-group peers.

Eliminates need for BGP.

MSDP Mesh-Groups

Configured with:

ip msdp mesh-group <name> <peer-address>

- Peers in the mesh-group must be fully meshed.
- Multiple mesh-groups per router are supported.

MSDP Mesh-Group Example



MSDP mesh-group peering

Avoid Mesh-Group Loops!!!

WARNING: There is no RPF check between Mesh-groups!!!



MSDP Mroute Flags

New 'mroute' Flags for MSDP



"A" flag indicates source is a candidate for advertisement by M

MSDP Enhancements

New IOS command

ip msdp new-rpf-rules

MSDP SA RPF check using IGP

Accept SA's from BGP NEXT HOP

Accept SA's from closest peer along the best path to the originating RP

"show ip msdp rpf"

MSDP RPF check using IGP

When MSDP peer = IGP peer (No BGP)

Find best IGP route to RP

Search URIB

If route to Originating RP found and:

If IGP next hop (or advertiser) address for RP is the

MSDP peer and in UP state, then that is the RPF

peer.

If route not found: Fall through to the next rule.

IGP Rule: MSDP peer = IGP peer (Next hop)



MSDP Peer = 3.1.1.1

IGP next hop to originating RP = 3.1.1.1

IGP next hop to originating RP = MSDP peer

SA RPF Check Succeeds



MSDP Peer

SA Message

IGP Rule: **MSDP** peer = **IGP** peer (Advertiser)



MSDP Peer = 4.1.1.1

IGP next hop to originating RP = 5.1.1.1

IGP advertiser to originating RP = 4.1.1.1

IGP advertiser to originating RP = MSDP peer

SA RPF Check Succeeds



SA's accepted from Next Hop


Accept SA along RPF path



Existing Rule: If first AS in best path to the RP != MSDP peer RPF Fails

New code: Choose peer in CLOSEST AS along best AS path to th Loosens rule a bit.



New MSDP RPF command

Router-A# show ip msdp rpf 2.1.1.1
RPF peer information for Router-B (2.1.1.1)
 RPF peer: Router-C (3.1.1.1)
 RPF route/mask: 2.1.1.0/24
 RPF rule: Peer is IGP next hop of best route
 RPF type: unicast (ospf 1)

Anycast-RP

- draft-ietf-mboned-anycast-rp-08.txt
- Within a domain, deploy more than one RP for the same group range
- Sources from one RP are known to other RPs using MSDP
- Give each RP the same /32 IP address
- Sources and receivers use closest RP, as determined by the IGP
- Used intra-domain to provide redundancy and RP load sharing, when an RP goes down, sources and receivers are taken to new RP via unicast routing



TECRST-1008_c1

Anycast-RP



Anycast-RP



MSDP Configuration



MSDP wrt SSM – Unnecessary!



MSDP wrt SSM – Unnecessary!



Agenda

- Introduction
- Multicast addressing
- Group Membership Protocol
- PIM-SM / SSM
- MSDP
- MBGP
- Summary

MBGP—Multiprotocol BGP

- MBGP overview
- MBGP capability negotiation
- MBGP NLRI exchange
- Configuration guidelines

MBGP

- Multiprotocol Extensions to BGP (RFC 2283).
- Tag unicast prefixes as multicast source prefixes for intradomain mcast routing protocols to do RPF checks.
- WHY? Allows for interdomain RPF checking where unicast and multicast paths are non-congruent.
- DO I REALLY NEED IT?
 - YES, if:

ISP to ISP peering

Multiple-homed networks

NO, if:

You are single-homed

MBGP Overview

 MBGP: Multiprotocol BGP (aka multicast BGP in multicast networks)

Defined in RFC 2283 (extensions to BGP)

Can carry different route types for different purposes

Unicast

Multicast

Both route types carried in same BGP session

Does not propagate multicast state information

Same path selection and validation rules

AS-Path, LocalPref, MED, ...

MBGP Overview

New multiprotocol attributes

MP_REACH_NLRI

MP_UNREACH_NLRI

MP_REACH_NLRI and MP_UNREACH_NLRI

Address Family Information (AFI) = 1 (IPv4)

Sub-AFI = 1 (NLRI is used for unicast)

Sub-AFI = 2 (NLRI is used for multicast RPF check)

Sub-AFI = 3 (NLRI is used for both unicast and multicast RPF check)

Allows for different policies between multicast and unicast

MBGP—Capability Negotiation

- BGP routers establish BGP sessions through the OPEN message
- OPEN message contains optional parameters
- BGP session is terminated if OPEN parameters are not recognised
- New parameter: CAPABILITIES

Multiprotocol extension

Multiple routes for same destination

Configures router to negotiate either or both NLRI

If neighbor configures both or subset, common NRLI is used in both directions

If there is no match, notification is sent and peering doesn't come up

If neighbor doesn't include the capability parameters in open, session backs off and reopens with no capability parameters

TECRST-1008_c1

302

MBGP—Summary

Solves part of inter-domain problem

Can exchange unicast prefixes for multicast RPF checks

Uses standard BGP configuration knobs

Permits separate unicast and multicast topologies if desired

Still must use PIM to:

Build distribution trees

Actually forward multicast traffic

PIM-SM recommended

MBGP configuration (new)

router bgp 1 address-family ipv4 unicast network 198.58.3.0/24 address-family ipv4 multicast network 198.58.3.0/24 neighbor 198.32.165.2 remote-as 2 description LabPeer1 update-source Ethernet0/0/1 address-family ipv4 unicast address-family ipv4 multicast

Your ASN

Configure prefixes to advertise in both SAFI-1 and SAFI-2

Your peer's ASN

Local address for the BGP peering session

Configure to exchange both SAFI-1 and SAFI-2 prefixes

MBGP configuration (original)



LAB #2 Interdomain Multicast

- Do not launch lab until instructred to do so.
- Lab templates or cfgs: Interdomain-Multicast
- Refer to your lab handout

LAB #2 Interdomain Multicast



#