

## **Deploying BGP**

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### **Presentation Slides**

### Slides are at:

#### ftp://ftp-eng.cisco.com /pfs/seminars/APRICOT2005-Deploying-BGP.pdf

#### And on the APRICOT2005 website

### Feel free to ask questions any time

### **BGP for Internet Service Providers**

- Scaling BGP
- Using Communities
- Deploying BGP in an ISP network

## **BGP Scaling Techniques**

### **BGP Scaling Techniques**

#### • How does a service provider:

Scale the iBGP mesh beyond a few peers?

Implement new policy without causing flaps and route churning?

Keep the network stable, scalable, as well as simple?

#### • Three Techniques:

Route Refresh, Flap Damping, Route Reflector

### **Route Refresh**

### **Route Refresh**

#### **Problem:**

- Hard BGP peer reset required after every policy change because the router does not store prefixes that are rejected by policy
- Hard BGP peer reset:

**Tears down BGP peering** 

**Consumes CPU** 

Severely disrupts connectivity for all networks

Solution:

Route Refresh

- Facilitates non-disruptive policy changes
- For most implementations, no configuration is needed

Automatically negotiated at peer establishment

- No additional memory is used
- Requires peering routers to support "route refresh capability" – RFC2918

### **Dynamic Reconfiguration**

- Use Route Refresh capability if supported find out from the BGP neighbour status display Non-disruptive, "Good For the Internet"
- If not supported, see if implementation has a workaround
- Only hard-reset a BGP peering as a last resort

#### Consider the impact to be equivalent to a router reboot

## **Route Flap Damping**

**Stabilising the Network** 

### **Route Flap Damping**

### Route flap

Going up and down of path or change in attribute BGP WITHDRAW followed by UPDATE = 1 flap eBGP neighbour peering reset is NOT a flap Ripples through the entire Internet Wastes CPU

 Damping aims to reduce scope of route flap propagation

### Route Flap Damping (continued)

#### Requirements

Fast convergence for normal route changes

**History predicts future behaviour** 

**Suppress oscillating routes** 

**Advertise stable routes** 

Documented in RFC2439

### Operation

# Add penalty for each flap NB: Change in attribute can also be penalized

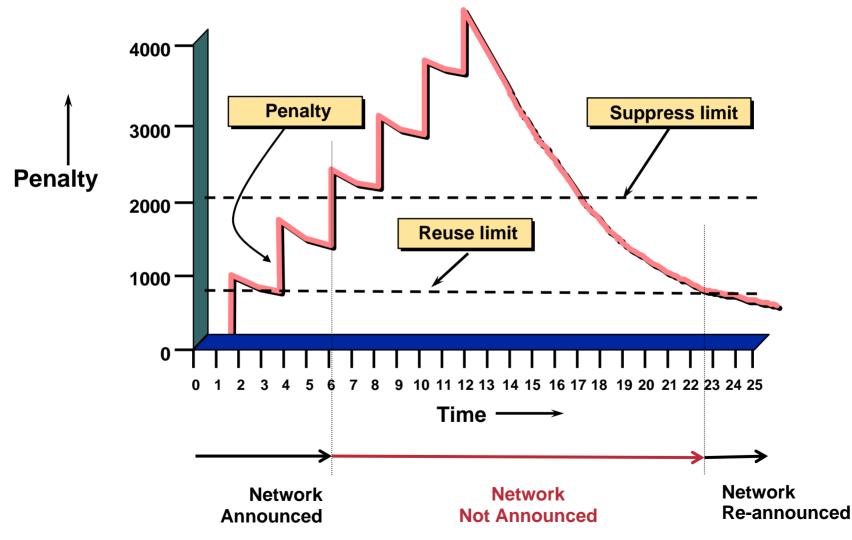
#### Exponentially decay penalty

half life determines decay rate

#### Penalty above suppress-limit do not advertise route to BGP peers

• Penalty decayed below reuse-limit re-advertise route to BGP peers

### Operation





- Only applied to inbound announcements from eBGP peers
- Alternate paths still usable
- Controllable by at least:
  - Half-life
  - reuse-limit
  - suppress-limit
  - maximum suppress time

### Configuration

#### Implementations allow various policy control with flap damping

Fixed damping, same rate applied to all prefixes

Variable damping, different rates applied to different ranges of prefixes

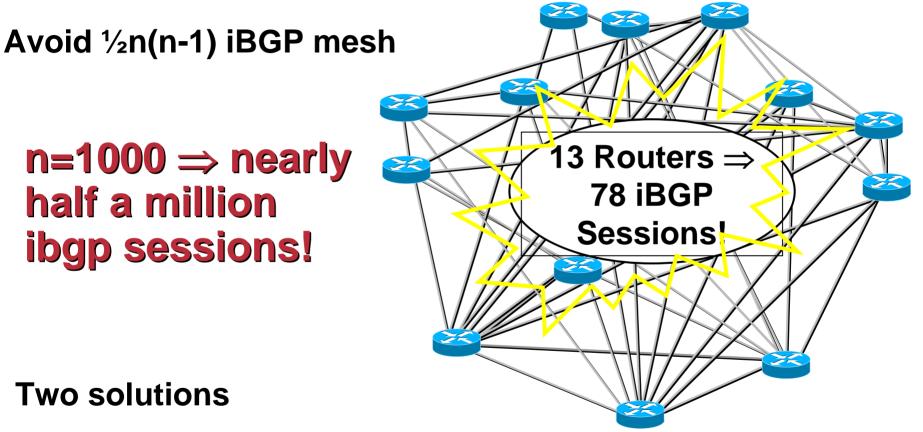
#### Recommendations for ISPs

http://www.ripe.net/docs/ripe-229.html

(work by European and US ISPs a few years ago as vendor defaults were considered to be too aggressive)

### **Route Reflectors**

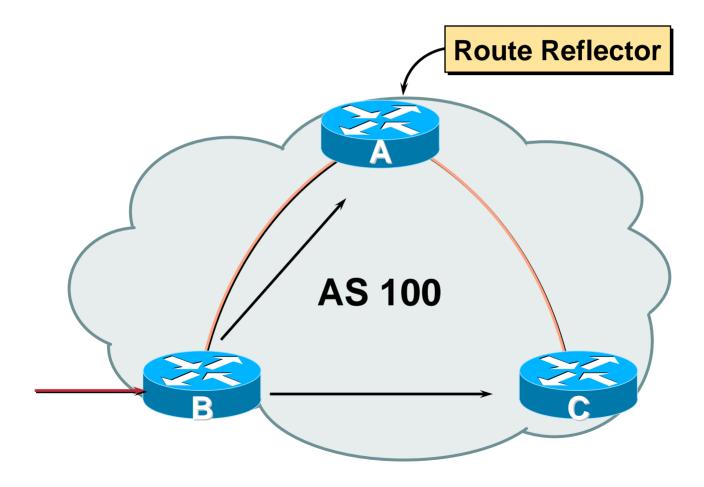
### Scaling iBGP mesh



Route reflector – simpler to deploy and run

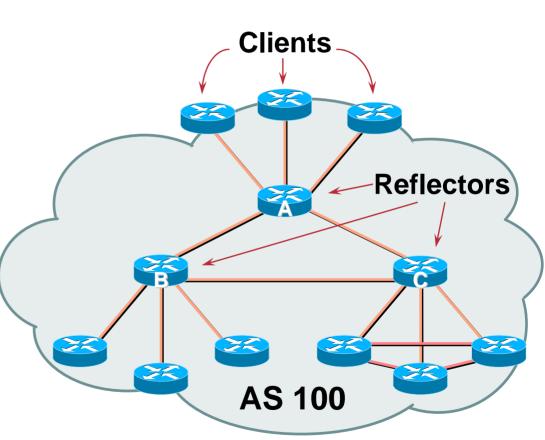
**Confederation – more complex, corner case benefits** 

### **Route Reflector: Principle**



### **Route Reflector**

- Reflector receives path from clients and nonclients
- Selects best path
- If best path is from client, reflect to other clients and non-clients
- If best path is from non-client, reflect to clients only
- Non-meshed clients
- Described in RFC2796



### **Route Reflector Topology**

- Divide the backbone into multiple clusters
- At least one route reflector and few clients per cluster
- Route reflectors are fully meshed
- Clients in a cluster could be fully meshed
- Single IGP to carry next hop and local routes

#### Route Reflectors: Loop Avoidance

#### Originator\_ID attribute

Carries the RID of the originator of the route in the local AS (created by the RR)

#### Cluster\_list attribute

The local cluster-id is added when the update is sent by the RR

Best to set cluster-id is from router-id (address of loopback)

(Some ISPs use their own cluster-id assignment strategy – but needs to be well documented!)

 Multiple RRs can be configured in the same cluster – not advised!

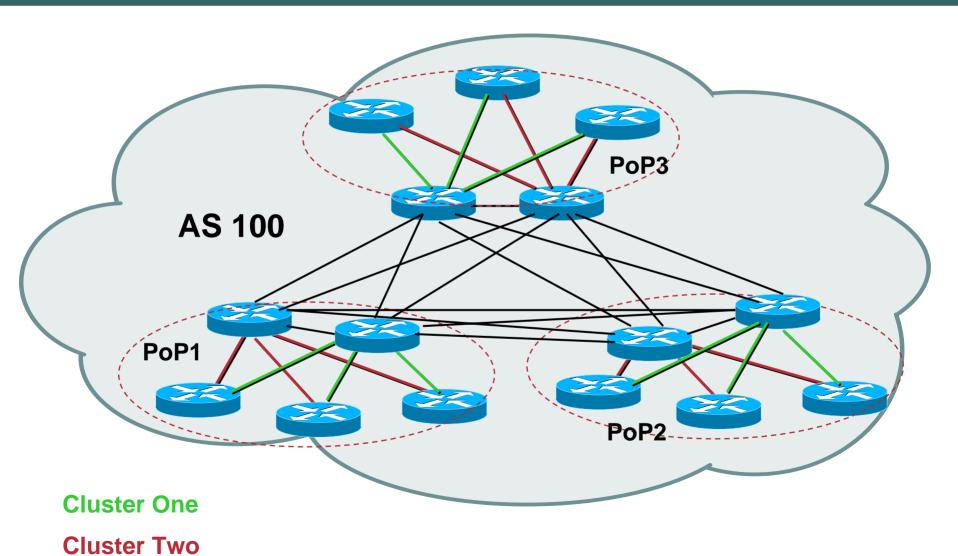
All RRs in the cluster must have the same cluster-id (otherwise it is a different cluster)

 A router may be a client of RRs in different clusters

Common today in ISP networks to overlay two clusters – redundancy achieved that way

 $\rightarrow$  Each client has two RRs = redundancy

#### Route Reflectors: Redundancy



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• Where to place the route reflectors? Always follow the physical topology!

This will guarantee that the packet forwarding won't be affected

• Typical ISP network:

**PoP has two core routers** 

**Core routers are RR for the PoP** 

**Two overlaid clusters** 

#### Route Reflectors: Migration

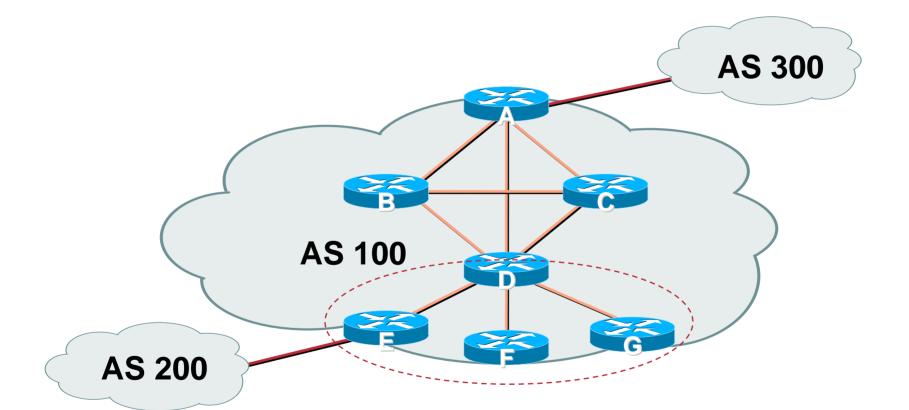
• Typical ISP network:

**Core routers have fully meshed iBGP** 

Create further hierarchy if core mesh too big Split backbone into regions

 Configure one cluster pair at a time Eliminate redundant iBGP sessions
 Place maximum one RR per cluster
 Easy migration, multiple levels

#### Route Reflector: Migration



#### Migrate small parts of the network, one part at a time.

### **BGP for Internet Service Providers**

- Scaling BGP
- Using Communities
- Deploying BGP in an ISP network

## Service Providers use of Communities

Some examples of how ISPs make life easier for themselves

- Another ISP "scaling technique"
- Prefixes are grouped into different "classes" or communities within the ISP network
- Each community means a different thing, has a different result in the ISP network

### **BGP Communities**

 Communities are generally set at the edge of the ISP network

**Customer edge:** customer prefixes belong to different communities depending on the services they have purchased

Internet edge: transit provider prefixes belong to difference communities, depending on the loadsharing or traffic engineering requirements of the local ISP, or what the demands from its BGP customers might be

 Two simple examples follow to explain the concept

- This demonstrates how communities might be used at the customer edge of an ISP network
- ISP has three connections to the Internet:

IXP connection, for local peers

Private peering with a competing ISP in the region

Transit provider, who provides visibility to the entire Internet

 Customers have the option of purchasing combinations of the above connections

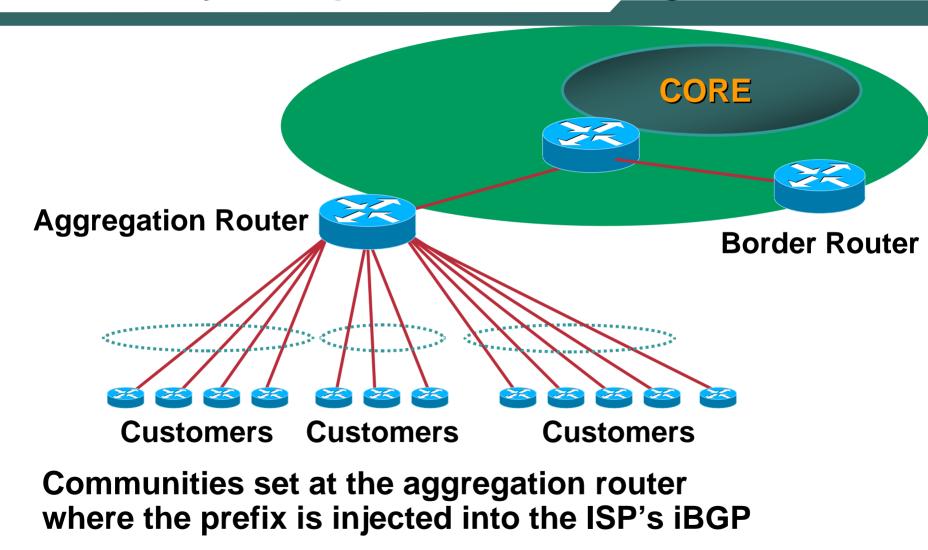
### **Community Example – Customer Edge**

• Community assignments:

IXP connection:community 100:2100Private peer:community 100:2200

- Customer who buys local connectivity (via IXP) is put in community 100:2100
- Customer who buys peer connectivity is put in community 100:2200
- Customer who wants both IXP and peer connectivity is put in 100:2100 and 100:2200
- Customer who wants "the Internet" has no community set
   We are going to announce his prefix everywhere

#### **Community Example – Customer Edge**



- No need to alter filters at the network border when adding a new customer
- New customer simply is added to the appropriate community

Border filters already in place take care of announcements

 $\Rightarrow$  Ease of operation!

### **Community Example – Internet Edge**

- This demonstrates how communities might be used at the peering edge of an ISP network
- ISP has four types of BGP peers:
  - Customer
  - **IXP** peer
  - **Private peer**
  - **Transit provider**
- The prefixes received from each can be classified using communities
- Customers can opt to receive any or all of the above

#### **Community Example – Internet Edge**

• Community assignments:

Customer prefix:	community 100:3000
IXP prefix:	community 100:3100
Private peer prefix:	community 100:3200

- BGP customer who buys local connectivity gets 100:3000
- BGP customer who buys local and IXP connectivity receives community 100:3000 and 100:3100
- BGP customer who buys full peer connectivity receives community 100:3000, 100:3100, and 100:3200
- Customer who wants "the Internet" gets everything Gets default route originated by aggregation router Or pays money to get all 135k prefixes

#### **Community Example – Internet Edge**

#### No need to create customised filters when adding customers

**Border router already sets communities** 

Installation engineers pick the appropriate community set when establishing the customer BGP session

 $\Rightarrow$  Ease of operation!

#### **Community Example – Summary**

- Two examples of customer edge and internet edge can be combined to form a simple community solution for ISP prefix policy control
- More experienced operators tend to have more sophisticated options available

Advice is to start with the easy examples given, and then proceed onwards as experience is gained

#### **Some ISP Examples**

- ISPs also create communities to give customers bigger routing policy control
- Public policy is usually listed in the IRR

Following examples are all in the IRR

Examples build on the configuration concepts from the introductory example

 Consider creating communities to give policy control to customers

**Reduces technical support burden** 

Reduces the amount of router reconfiguration, and the chance of mistakes

#### **Some ISP Examples: Sprintlink**



http://www.sprintlink.net/policy/bgp.html

#### WHAT YOU CAN CONTROL

AS-PATH PREPENDS

Sprint allows customers to use AS-path prepending to adjust route preference on the network. Such prepending will be received and passed on properly without notifiving Sprint of your change in announcments.

Additionally, Sprint will prepend AS1239 to eBGP sessions with certain autonomous systems depending on a received community. Currently, the following ASes are supported: 1668, 209, 2914, 3300, 3356, 3549, 3561, 4635, 701, 7018, 702 and 8220.

String	Resulting AS Path to ASXXX	
65000:XX	C Do not advertise to ASXXX	
65001:XX	< 1239 (default)	
65002:XX	< 1239 1239	
65003:XX	< 1239 1239 1239	
65004:XXX	1239 1239 1239 1239	
String	Resulting AS Path to ASXXX in Asia	
65070:XXX	Do not advertise to ASXXX	
65071:XXX	1239 (default)	
65072:XXX	1239 1239	
65073:XXX	1239 1239 1239	
65074:XXX	1239 1239 1239 1239	
String Resulting AS Path to ASXXX in Europe		
65050:XXX	Do not advertise to ASXXX	
65051:XXX	1239 (default)	
65052:XXX	1239 1239	
65053:XXX	1239 1239 1239	
65054:XXX	1239 1239 1239 1239	
04-1-1-1	Resulting AS Path to ASXXX in North	
String	America	
65010:XXX	Do not advertise to ASXXX	
65011:XXX	1239 (default)	
65012:XXX	1239 1239	
65013:XXX	1239 1239 1239	
65014:XXX	1239 1239 1239 1239	
String R	esulting AS Path to all supported ASes	
65000:0	Do not advertise	
65001:0	1239 (default)	
	1000	

1239 1239 ...

65002:0

More info at

www.sprintlink.net/policy/bgp.html

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#### Some ISP Examples Connect.com.au

aut-num:	AS2764
descr:	connect.com.au pty 1td
remarks:	Community Definition
remarks:	
remarks:	2764:1 Announce to "domestic" rate ASes only
remarks:	2764:2 Don't announce outside local POP
remarks:	2764:3 Lower local preference by 25
remarks:	2764:4 Lower local preference by 15
remarks:	2764:5 Lower local preference by 5
remarks:	2764:6 Announce to non customers with "no-export"
remarks:	2764:7 Only announce route to customers
remarks:	2764:8 Announce route over satellite link
notify:	routing@connect.com.au
mnt-by:	CONNECT-AU
changed:	mrp@connect.com.au 19990506
source:	CCAIR

#### More at http://info.connect.com.au/docs/routing/general/multi-faq.shtml#q13

#### Some ISP Examples MCI Europe

aut-num:	AS702		
descr:	MCI EMEA	- Commercial IP service provider in	Europe
remarks:	MCI uses	the following communities with its o	customers:
	702:80	Set Local Pref 80 within AS702	
	702:120	Set Local Pref 120 within AS702	
	702:20	Announce only to MCI AS'es and MCI	customers
	702:30	Keep within Europe, don't announce	
	702:1	Prepend AS702 once at edges of MCI	
	702:2	Prepend AS702 twice at edges of MCI	to Peers
	702:3	Prepend AS702 thrice at edges of MC	I to Peers
	Advanced communities for customers		
	702:7020	Do not announce to AS702 peers with	n a scope of
		National but advertise to Global Pe	eers, European
		Peers and MCI customers.	
	702:7001	Prepend AS702 once at edges of MCI	to AS702
		peers with a scope of National.	
<snip></snip>			
	Additional details of the MCI communities are located at:		
	http://global.mci.com/uk/customer/bgp/		
mnt-by:	WCOM-EMEA-RICE-MNT And Several		
changed:	rice@lists.mci.com 20041006 more!		
source:	RIPE		

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#### Some ISP Examples BT Ignite

aut-num:	AS5400	
descr:	BT Ignite European Backbone	
remarks:		
remarks:	Community to	Community to
remarks:	Not announce To peer:	AS prepend 5400
remarks:		
remarks:	5400:1000 All peers & Trans:	its 5400:2000
remarks:		
remarks:	5400:1500 All Transits	5400:2500
remarks:	5400:1501 Sprint Transit (As	s1239) 5400:2501
remarks:	5400:1502 SAVVIS Transit (As	s3561) 5400:2502
remarks:	5400:1503 Level 3 Transit (A	AS3356) 5400:2503
remarks:	5400:1504 AT&T Transit (AS70	018) 5400:2504
remarks:	5400:1505 UUnet Transit (AS	701) 5400:2505
remarks:		
remarks:	5400:1001 Nexica (AS24592)	5400:2001
remarks:	5400:1002 Fujitsu (AS3324)	5400:2002
remarks:	5400:1003 Unisource (AS3300)	) 5400:2003
<snip></snip>		
notify:	notify@eu.bt.net	And many
mnt-by:	CIP-MNT	
source:	RIPE	many more!

#### Some ISP Examples Carrier1

aut-num:	AS8918
descr:	Carrier1 Autonomous System
<snip></snip>	
remarks:	Community Definition
remarks:	*
remarks:	8918:2000 Do not announce to C1 customers
remarks:	8918:2010 Do not announce to C1 peers, peers+ and transit
remarks:	8918:2015 Do not announce to C1 transit providers
remarks:	*
remarks:	8918:2020 Do not announce to Teleglobe (AS 6453)
remarks:	8918:2035 Do not announce to UUNet (AS 702)
remarks:	8918:2040 Do not announce to Cogent (AS 174)
remarks:	8918:2050 Do not announce to T-Systems (AS 3320)
remarks:	8918:2060 Do not announce to Sprint (AS 1239)
remarks:	*
remarks:	8918:2070 Do not announce to AMS-IX peers
remarks:	8918:2080 Do not announce to NL-IX peers
remarks:	8918:2090 Do not announce to Packet Exchange Peers
<snip></snip>	
notify:	inoc@carrier1.net And many
mnt-by:	CARRIER1-MNT many more!
source:	RIPE

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#### Some ISP Examples Level 3

aut-num:	AS3356
descr:	Level 3 Communications
<snip></snip>	
remarks:	
remarks:	customer traffic engineering communities - Suppression
remarks:	
remarks:	64960:XXX - announce to AS XXX if 65000:0
remarks:	65000:0 - announce to customers but not to peers
remarks:	65000:XXX - do not announce at peerings to AS XXX
remarks:	
remarks:	customer traffic engineering communities - Prepending
remarks:	
remarks:	65001:0 - prepend once to all peers
remarks:	65001:XXX - prepend once at peerings to AS XXX
remarks:	65002:0 - prepend twice to all peers
remarks:	65002:XXX - prepend twice at peerings to AS XXX
remarks:	65003:0 - prepend 3x to all peers
remarks:	65003:XXX - prepend 3x at peerings to AS XXX
remarks:	65004:0 - prepend 4x to all peers
remarks:	65004:XXX - prepend 4x at peerings to AS XXX
<snip></snip>	
mnt-by:	LEVEL3-MNT And many
source:	RIPE many more!

#### **BGP for Internet Service Providers**

- Scaling BGP
- Using Communities
- Deploying BGP in an ISP network

## Deploying BGP in an ISP Network

Okay, so we've learned all about BGP now; how do we use it on our network??

#### **Deploying BGP**

- The role of IGPs and iBGP
- Aggregation
- Receiving Prefixes
- Configuration Tips

# The role of IGP and iBGP

Ships in the night?

#### Or

**Good foundations?** 

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# Internal Routing Protocols (IGPs) examples are ISIS and OSPF

used for carrying infrastructure addresses

**NOT** used for carrying Internet prefixes or customer prefixes

design goal is to minimise number of prefixes in IGP to aid scalability and rapid convergence

#### BGP used internally (iBGP) and externally (eBGP)

#### iBGP used to carry

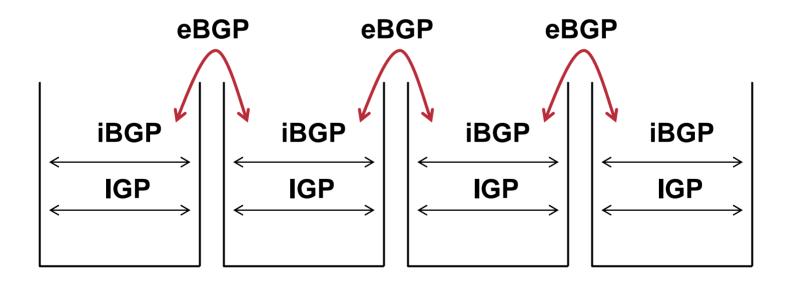
some/all Internet prefixes across backbone customer prefixes

#### eBGP used to

exchange prefixes with other ASes implement routing policy

#### **BGP/IGP model used in ISP networks**

#### Model representation



#### **BGP versus OSPF/ISIS**

#### • DO NOT:

# distribute BGP prefixes into an IGP distribute IGP routes into BGP use an IGP to carry customer prefixes YOUR NETWORK WILL NOT SCALE

#### Injecting prefixes into iBGP

- Use iBGP to carry customer prefixes don't ever use IGP
- Point static route to customer interface
- Enter network into BGP process

Ensure that implementation options are used so that the prefix always remains in iBGP, regardless of state of interface

i.e. avoid iBGP flaps caused by interface flaps

# Aggregation

**Quality or Quantity?** 



- Aggregation means announcing the address block received from the RIR to the other ASes connected to your network
- Subprefixes of this aggregate may be:

Used internally in the ISP network

Announced to other ASes to aid with multihoming

 Unfortunately too many people are still thinking about class Cs, resulting in a proliferation of /24s in the Internet routing table



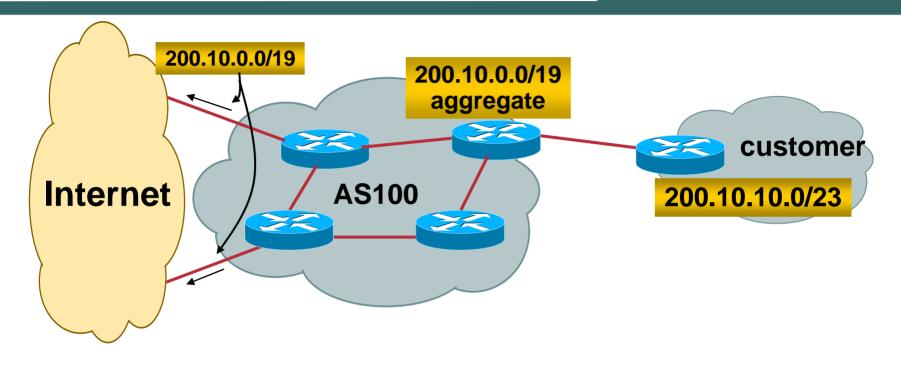
- Address block should be announced to the Internet as an aggregate
- Subprefixes of address block should NOT be announced to Internet unless special circumstances (more later)
- Aggregate should be generated internally Not on the network borders!

#### **Announcing an Aggregate**

- ISPs who don't and won't aggregate are held in poor regard by community
- Registries publish their minimum allocation size Anything from a /20 to a /22 depending on RIR Different sizes for different address blocks
- No real reason to see anything longer than a /22 prefix in the Internet

BUT there are currently >84000 /24s!

#### **Aggregation – Example**



- Customer has /23 network assigned from AS100's /19 address block
- AS100 announced /19 aggregate to the Internet

#### **Aggregation – Good Example**

Customer link goes down

their /23 network becomes unreachable

/23 is withdrawn from AS100's iBGP

 /19 aggregate is still being announced

no BGP hold down problems

no BGP propagation delays

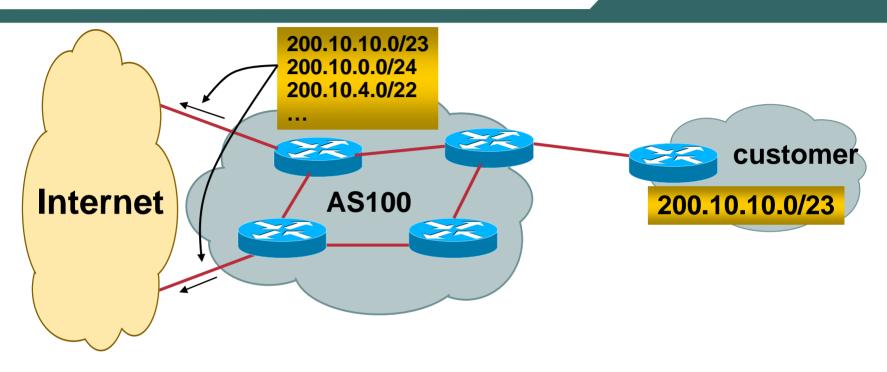
no damping by other ISPs

- Customer link returns
  - Their /23 network is visible again

The /23 is re-injected into AS100's iBGP

- The whole Internet becomes visible immediately
- Customer has Quality of Service perception

#### **Aggregation – Example**



- Customer has /23 network assigned from AS100's /19 address block
- AS100 announces customers' individual networks to the Internet

#### **Aggregation – Bad Example**

#### Customer link goes down

Their /23 network becomes unreachable

/23 is withdrawn from AS100's iBGP

 Their ISP doesn't aggregate its /19 network block

/23 network withdrawal announced to peers

starts rippling through the Internet

added load on all Internet backbone routers as network is removed from routing table

#### → • Customer link returns

Their /23 network is now visible to their ISP

Their /23 network is readvertised to peers

Starts rippling through Internet

Load on Internet backbone routers as network is reinserted into routing table

Some ISP's suppress the flaps

Internet may take 10-20 min or longer to be visible

Where is the Quality of Service???

# Good example is what everyone should do! Adds to Internet stability Reduces size of routing table Reduces routing churn Improves Internet QoS for everyone Bad example is what too many still do!

Why? Lack of knowledge? Laziness?

#### The Internet Today (February 2005)

 Current Internet Routing Table Statistics **BGP Routing Table Entries** 154984 **Prefixes after maximum aggregation** 90381 **Unique prefixes in Internet** 74096 **Prefixes smaller than registry alloc** 72278 **/24s announced** 84524 only 5646 /24s are from 192.0.0/8 ASes in use 18880

#### Efforts to improve aggregation

#### • The CIDR Report

Initiated and operated for many years by Tony Bates

Now combined with Geoff Huston's routing analysis

www.cidr-report.org

Results e-mailed on a weekly basis to most operations lists around the world

Lists the top 30 service providers who could do better at aggregating

#### Efforts to improve aggregation The CIDR Report

- Also computes the size of the routing table assuming ISPs performed optimal aggregation
- Website allows searches and computations of aggregation to be made on a per AS basis

flexible and powerful tool to aid ISPs

Intended to show how greater efficiency in terms of BGP table size can be obtained without loss of routing and policy information

Shows what forms of origin AS aggregation could be performed and the potential benefit of such actions to the total table size

Very effectively challenges the traffic engineering excuse



#### Aggregation on the Internet could be MUCH better

35% saving on Internet routing table size is quite feasible

**Tools are available** 

Commands on the router are not hard

**CIDR-Report webpage** 

# **Receiving Prefixes**

#### **Receiving Prefixes**

 There are three scenarios for receiving prefixes from other ASNs

**Customer talking BGP** 

**Peer talking BGP** 

**Upstream/Transit talking BGP** 

 Each has different filtering requirements and need to be considered separately

#### **Receiving Prefixes: From Customers**

- ISPs should only accept prefixes which have been assigned or allocated to their downstream customer
- If ISP has assigned address space to its customer, then the customer IS entitled to announce it back to his ISP
- If the ISP has NOT assigned address space to its customer, then:

Check in the four RIR databases to see if this address space really has been assigned to the customer

The tool: whois –h whois.apnic.net x.x.x.0/24

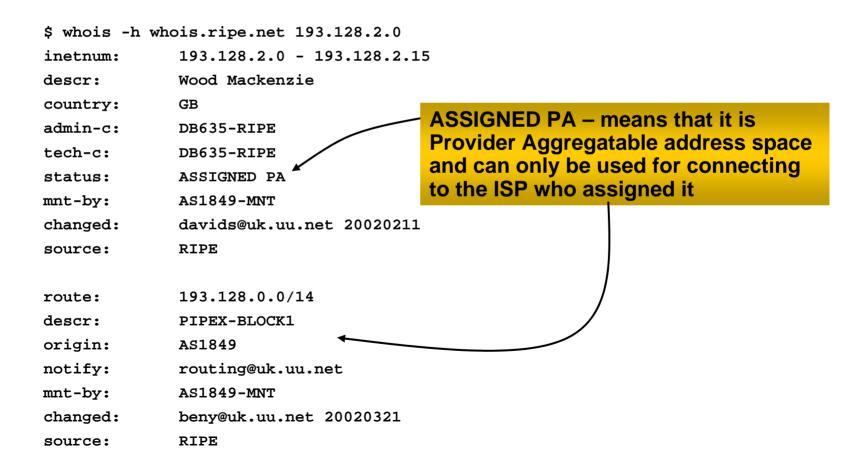
#### **Receiving Prefixes: From Customers**

 Example use of whois to check if customer is entitled to announce address space:

pfs-pc\$ whois	-h whois.apnic.n	et 202.12.29.0
inetnum:	202.12.29.0 - 20	2.12.29.255
netname:	APNIC-AP-AU-BNE	
descr:	APNIC Pty Ltd - Brisbane Offices + Servers	
descr:	Level 1, 33 Park Rd	
descr:	PO Box 2131, Milton	
descr:	Brisbane, QLD.	
country:	AU	Portable – means its an assignment
admin-c:	HM20-AP	to the customer, the customer can
tech-c:	NO4-AP	announce it to you
mnt-by:	APNIC-HM	
-1 1	hm-changed@apnic.net 20030108	
changed:	hm-changed@apnic	.net 20030108 /
changed: status:	hm-changed@apnic ASSIGNED PORTABL	

#### **Receiving Prefixes: From Customers**

 Example use of whois to check if customer is entitled to announce address space:



 A peer is an ISP with whom you agree to exchange prefixes you originate into the Internet routing table

Prefixes you accept from a peer are only those they have indicated they will announce

Prefixes you announce to your peer are only those you have indicated you will announce

#### **Receiving Prefixes: From Peers**

- Agreeing what each will announce to the other:
  - Exchange of e-mail documentation as part of the peering agreement, and then ongoing updates

#### OR

Use of the Internet Routing Registry and configuration tools such as the IRRToolSet

http://www.isc.org/sw/IRRToolSet/

#### **Receiving Prefixes: From Upstream/Transit Provider**

- Upstream/Transit Provider is an ISP who you pay to give you transit to the WHOLE Internet
- Receiving prefixes from them is not desirable unless really necessary

special circumstances – see Multihoming Tutorial

 Ask upstream/transit provider to either: originate a default-route OR

announce one prefix you can use as default

#### **Receiving Prefixes: From Upstream/Transit Provider**

 If necessary to receive prefixes from any provider, care is required

don't accept RFC1918 etc prefixes

ftp://ftp.rfc-editor.org/in-notes/rfc3330.txt

don't accept your own prefixes

don't accept default (unless you need it)

don't accept prefixes longer than /24

 Check Rob Thomas' list of "bogons" http://www.cymru.com/Documents/bogon-list.html  Paying attention to prefixes received from customers, peers and transit providers assists with:

The integrity of the local network

The integrity of the Internet

 Responsibility of all ISPs to be good Internet citizens

### **Preparing the Network**

- We want to deploy BGP now...
- BGP will be used therefore an ASN is required
- If multihoming to different ISPs is intended in the near future, a public ASN should be obtained:

Either go to upstream ISP who is a registry member, or

Apply to the RIR yourself for a one off assignment, or

Ask an ISP who is a registry member, or

Join the RIR and get your own IP address allocation too (this option strongly recommended)!

 The network is not running any BGP at the moment

single statically routed connection to upstream ISP

 The network is not running any IGP at all Static default and routes through the network to do "routing"

- Decide on IGP: OSPF or ISIS ③
- Assign loopback interfaces and /32 addresses to each router which will run the IGP

Loopback is used for OSPF and BGP router id anchor Used for iBGP and route origination

#### Deploy IGP (e.g. OSPF)

IGP can be deployed with NO IMPACT on the existing static routing

e.g. OSPF distance might be 110, static distance is 1

**Smallest distance wins** 

 Be prudent deploying IGP – keep the Link State Database Lean!

**Router loopbacks go in IGP** 

Backbone WAN point to point links go in IGP

(In fact, any link where IGP dynamic routing will be run should go into IGP)

Summarise on area/level boundaries (if possible) – i.e. think about your IGP address plan

#### • Routes which don't go into the IGP include:

**Dynamic assignment pools (DSL/Cable/Dial)** 

**Customer point to point link addressing** 

(using next-hop-self in iBGP ensures that these do NOT need to be in IGP)

**Static/Hosting LANs** 

**Customer assigned address space** 

Anything else not listed in the previous slide

# Preparing the Network iBGP

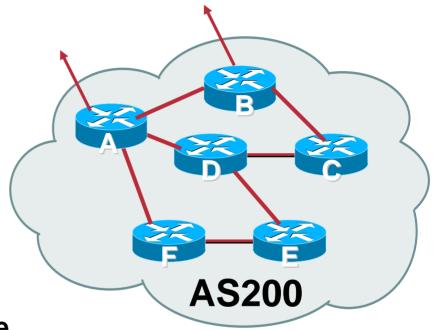
- Second step is to configure the local network to use iBGP
- iBGP can run on

all routers, or

a subset of routers, or

just on the upstream edge

• *iBGP must run on all routers which are in the transit path between external connections* 



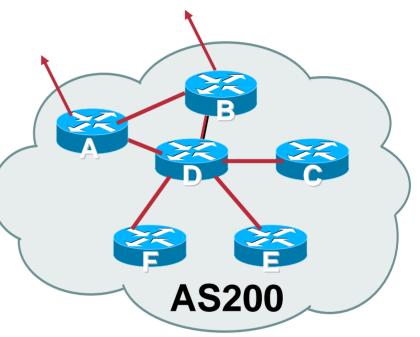
#### Preparing the Network iBGP (Transit Path)

- iBGP must run on all routers which are in the transit path between external connections
- Routers C, E and F are not in the transit path

Static routes or IGP will suffice

 Router D is in the transit path

> Will need to be in iBGP mesh, otherwise routing loops will result



# Typical SP networks have three layers: Core – the backbone, usually the transit path Distribution – the middle, PoP aggregation layer

Aggregation – the edge, the devices connecting customers

#### Preparing the Network Aggregation Layer

#### iBGP is optional

Many ISPs run iBGP here, either partial routing (more common) or full routing (less common)

Full routing is not needed unless customers want full table

Partial routing is cheaper/easier, might usually consist of internal prefixes and, optionally, external prefixes to aid external load balancing

Communities make this administratively easy

Many aggregation devices can't run iBGP
 Static routes from distribution devices for address pools
 IGP for best exit

#### **Preparing the Network Distribution Layer**

#### Usually runs iBGP

Partial or full routing (as with aggregation layer)

#### But does not have to run iBGP

IGP is then used to carry customer prefixes (does not scale)

**IGP** is used to determine nearest exit

#### Networks which plan to grow large should deploy iBGP from day one

Migration at a later date is extra work

No extra overhead in deploying iBGP; indeed, the IGP benefits

- Core of network is usually the transit path
- iBGP necessary between core devices
   Full routes or partial routes:
   Transit ISPs carry full routes in core
   Edge ISPs carry partial routes only

   Core layer includes AS border routers

Decide on:

Best iBGP policy

Will it be full routes everywhere, or partial, or some mix?

• iBGP scaling technique

**Community policy?** 

**Route-reflectors?** 

**Techniques such as peer templates?** 

#### Then deploy iBGP:

**Step 1: Introduce iBGP mesh on chosen routers** 

make sure that iBGP distance is greater than IGP distance (it usually is)

Step 2: Install "customer" prefixes into iBGP

**Check!** Does the network still work?

Step 3: Carefully remove the static routing for the prefixes now in IGP and iBGP

**Check!** Does the network still work?

**Step 4: Deployment of eBGP follows** 

#### Install "customer" prefixes into iBGP?

• Customer assigned address space

**Network statement/static route combination** 

Use unique community to identify customer assignments

#### Customer facing point-to-point links

Redistribute connected routes through filters which only permit point-to-point link addresses to enter iBGP

Use a unique community to identify point-to-point link addresses (these are only required for your monitoring system)

#### Dynamic assignment pools & local LANs

Simple network statement will do this

Use unique community to identify these networks

#### Carefully remove static routes?

• Work on one router at a time:

Check that static route for a particular destination is also learned either by IGP or by iBGP

If so, remove it

If not, establish why and fix the problem

(Remember to look in the RIB, not the FIB!)

- Then the next router, until the whole PoP is done
- Then the next PoP, and so on until the network is now dependent on the IGP and iBGP you have deployed

#### Preparing the Network Completion

#### Previous steps are NOT flag day steps

Each can be carried out during different maintenance periods, for example:

**Step One on Week One** 

**Step Two on Week Two** 

**Step Three on Week Three** 

And so on

And with proper planning will have NO customer visible impact at all

#### Preparing the Network Configuration Summary

- IGP essential networks are in IGP
- Customer networks are now in iBGP iBGP deployed over the backbone
   Full or Partial or Upstream Edge only
- BGP distance is greater than any IGP
- Now ready to deploy eBGP

## **Configuration Tips**

Of templates, passwords, tricks, and more templates

- Make sure loopback is configured on router iBGP between loopbacks, NOT real interfaces
- Make sure IGP carries loopback /32 address
- Consider the DMZ nets:

Use unnumbered interfaces? Use next-hop-self on iBGP neighbours Or carry the DMZ /30s in the iBGP Basically keep the DMZ nets out of the IGP!

#### **Next-hop-self**

Used by many ISPs on edge routers

Preferable to carrying DMZ /30 addresses in the IGP

**Reduces size of IGP to just core infrastructure** 

Alternative to using unnumbered interfaces

Helps scale network

BGP speaker announces external network using local address (loopback) as next-hop

#### **Templates**

 Good practice to configure templates for everything

Vendor defaults tend not to be optimal or even very useful for ISPs

ISPs create their own defaults by using configuration templates

eBGP and iBGP examples follow

Also see Project Cymru's BGP templates

www.cymru.com/Documents

#### **iBGP Template** Example

- iBGP between loopbacks!
- Next-hop-self

Keep DMZ and external point-to-point out of IGP

- Always send communities in iBGP
   Otherwise accidents will happen
- Hardwire BGP to version 4

Yes, this is being paranoid!

Use passwords on iBGP session
 Not being paranoid, VERY necessary

#### eBGP Template Example

BGP damping

Use RIPE-229 parameters, or something even weaker Don't use the vendor defaults without thinking

Remove private ASes from announcements

**Common omission today** 

Use extensive filters, with "backup"

Use as-path filters to backup prefix filters

Keep policy language for implementing policy, rather than basic filtering

Use password agreed between you and peer on eBGP session

#### eBGP Template Example continued

Use maximum-prefix tracking

Router will warn you if there are sudden increases in BGP table size, bringing down eBGP if desired

#### Log changes of neighbour state

...and monitor those logs!

 Make BGP admin distance higher than that of any IGP

Otherwise prefixes heard from outside your network could override your IGP!!

#### Limiting AS Path Length

 Some BGP implementations have problems with long AS\_PATHS

**Memory corruption** 

**Memory fragmentation** 

 Even using AS\_PATH prepends, it is not normal to see more than 20 ASes in a typical AS\_PATH in the Internet today

The Internet is around 5 ASes deep on average Largest AS\_PATH is usually 16-20 ASNs

#### Some announcements have ridiculous lengths of AS-paths:

\*> 3FFE:1600::/24 3FFE:C00:8023:5::2 22 11537 145 12199
10318 10566 13193 1930 2200 3425 293 5609 5430 13285 6939
14277 1849 33 15589 25336 6830 8002 2042 7610 i

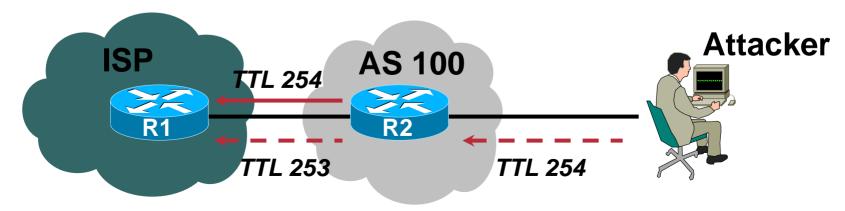
This example is an error in one IPv6 implementation

 If your implementation supports it, consider limiting the maximum AS-path length you will accept • Implement RFC3682 on BGP peerings

Neighbour sets TTL to 255

Local router expects TTL of incoming BGP packets to be 254

No one apart from directly attached devices can send BGP packets which arrive with TTL of 254, so any possible attack by a remote miscreant is dropped due to TTL mismatch



#### **BGP TTL "hack"**

#### • TTL Hack:

Both neighbours must agree to use the feature TTL check is much easier to perform than MD5 (Called BTSH – BGP TTL Security Hack)

#### Provides "security" for BGP sessions

In addition to packet filters of course

MD5 should still be used for messages which slip through the TTL hack

See www.nanog.org/mtg-0302/hack.html for more details

#### **Passwords on BGP sessions**

- Yes, I am mentioning passwords again
- Put password on the BGP session

It's a secret shared between you and your peer

If arriving packets don't have the correct MD5 hash, they are ignored

Helps defeat miscreants who wish to attack BGP sessions

 Powerful preventative tool, especially when combined with filters and the TTL "hack"



- Use configuration templates
- Standardise the configuration
- Be aware of standard "tricks" to avoid compromise of the BGP session
- Anything to make your life easier, network less prone to errors, network more likely to scale
- It's all about scaling if your network won't scale, then it won't be successful

#### **BGP for Internet Service Providers**

- Scaling BGP
- Using Communities
- Deploying BGP in an ISP network



# **Deploying BGP**

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